

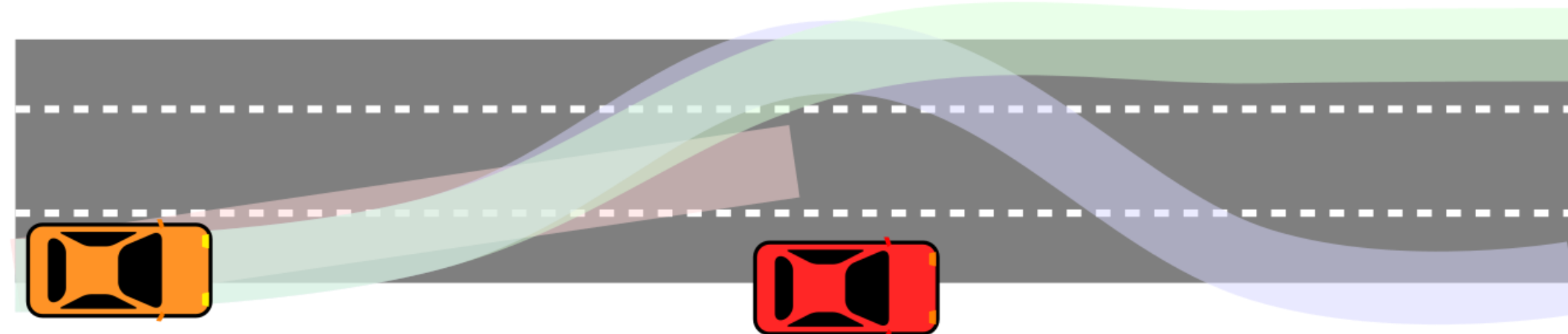
Autonomous Systems: Specification and Verification

Lancaster University
School of Computing and Communications

Andrew Sogokon (a.sogokon@lancaster.ac.uk)
Prof Neeraj Suri (PI)

Specifications for Autonomous Systems

- Specifications are descriptions of what a system should (or should not) do.
- A large source of specifications for AS comes from **regulations** (e.g. the *Highway Code* for terrestrial vehicles, or the *Rules of the Air* for aerial vehicles).
- Regulations written in natural language (e.g. English prose) can be imprecise and subject to various interpretations.
- E.g. "When changing the lane to the left lane during overtaking, no following road user shall be endangered" (Rizaldi et al., 2017).



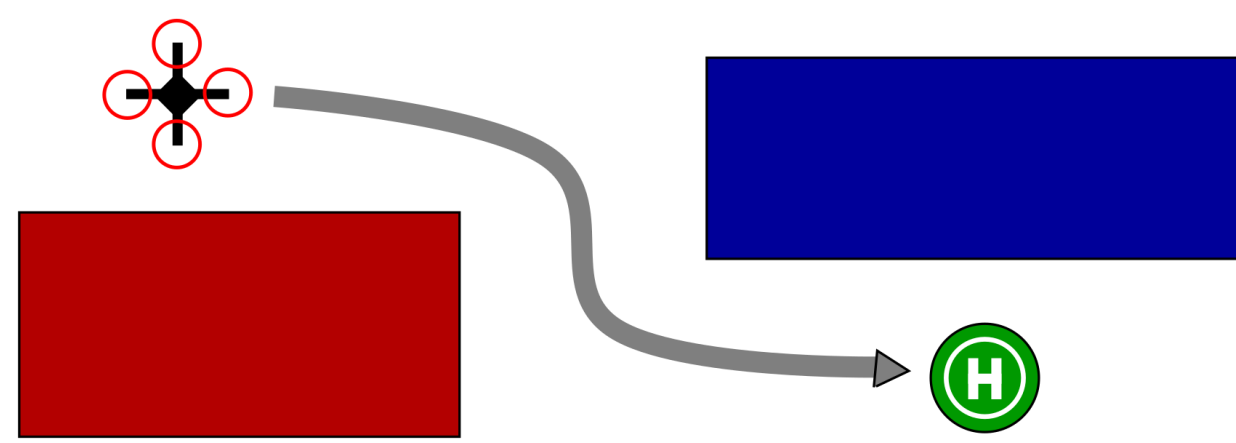
Formal Specifications

- A **formal model** of the system provides a precise description of the dynamics.
- A **formal specification** can be *verified* against a formal model.
- Mission specifications can often be stated in formal logic (such as various **temporal logics**) and can incorporate safety and liveness requirements:

$$\neg \square_{[0, \infty)} \text{dist}(\text{ownership}, \text{intruder}) \geq d_{\min} \quad (\text{collision avoidance})$$



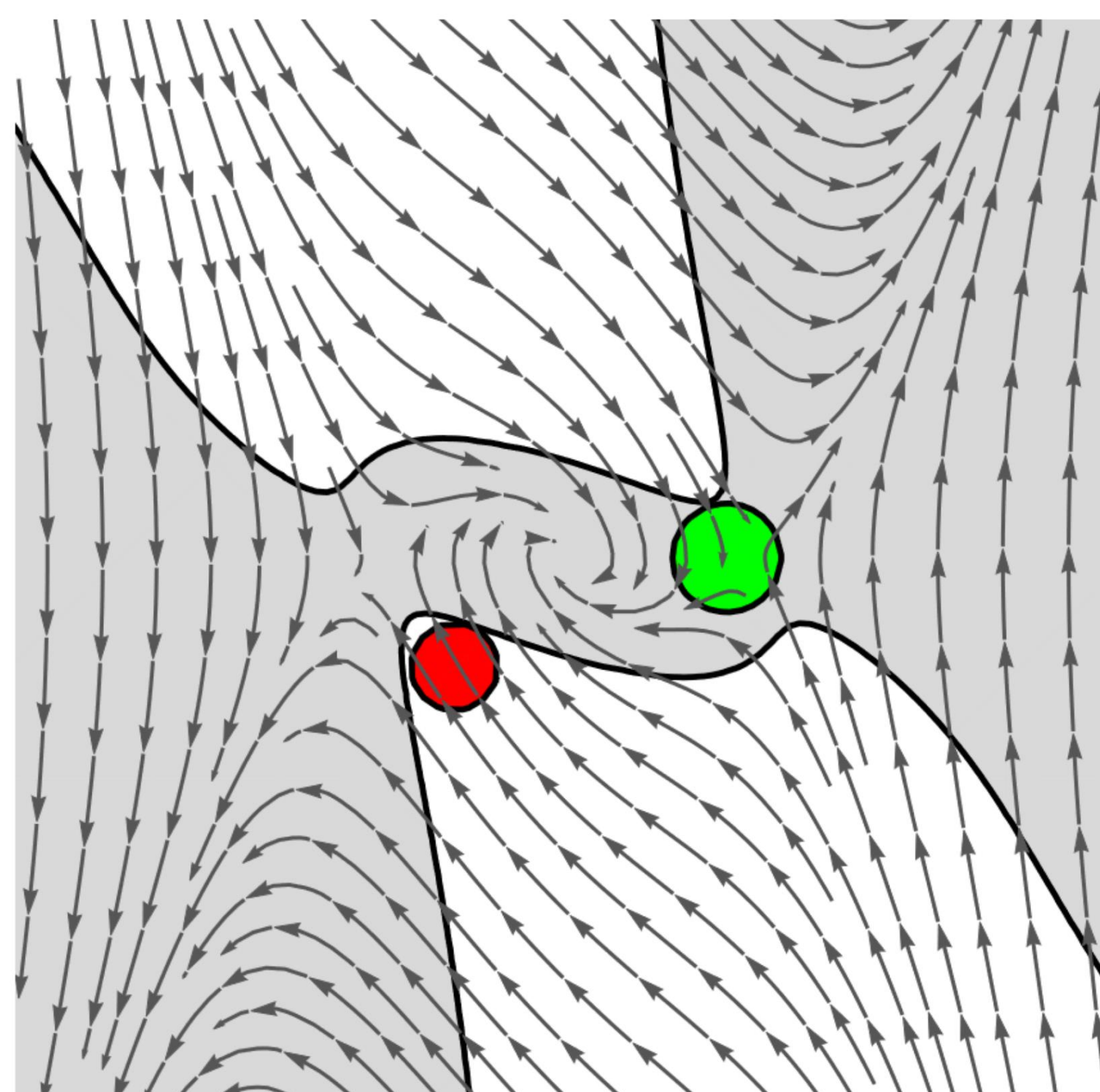
$$\neg (\diamond_{[0, T]} \text{Target}) \wedge (\square_{[0, T]} \text{Safe}) \quad (\text{reach-avoid})$$



Safety Specification and Verification

Safety Specifications

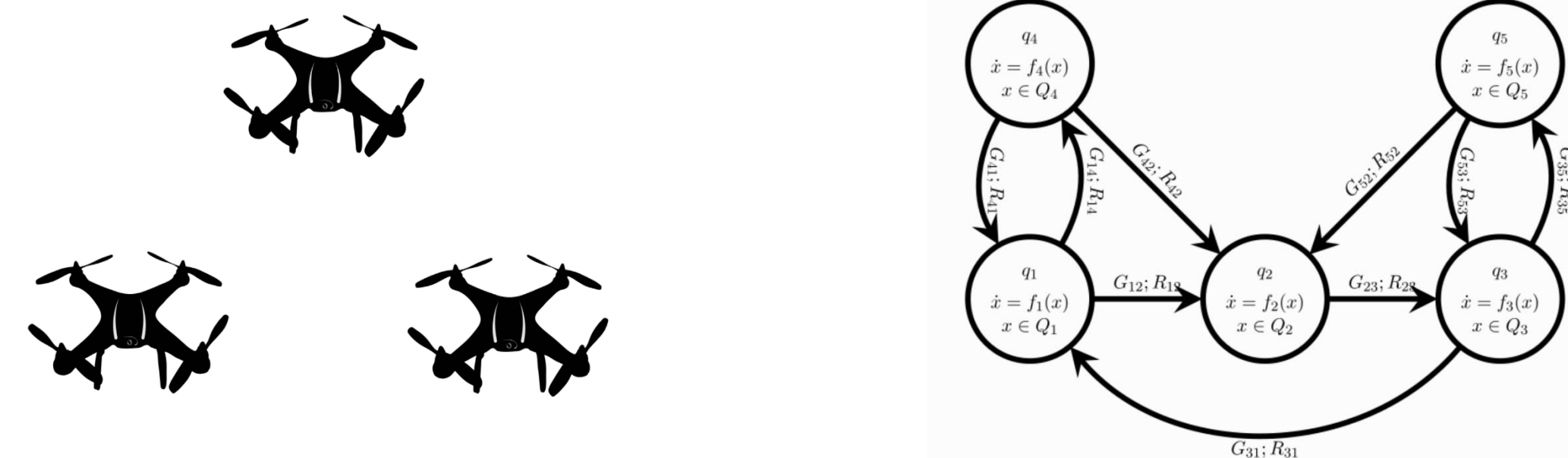
- A **safety specification** for a given system requires two elements:
 - 1 - A description of the possible initial states from which the system may begin its operation.
 - 2 - A description of undesirable (i.e. unsafe) states into which the system must never transition.



- Safety verification** is concerned with proving a safety specification, i.e. rigorously demonstrating that a system may never transition into any of the unsafe states provided that it starts operating from one of the specified initial states.

Cyber-Physical Autonomous Systems

- Systems that interact with a physical environment are *cyber-physical systems (CPS)*.
- Continuous dynamics in CPS is usually described using **differential equations**.



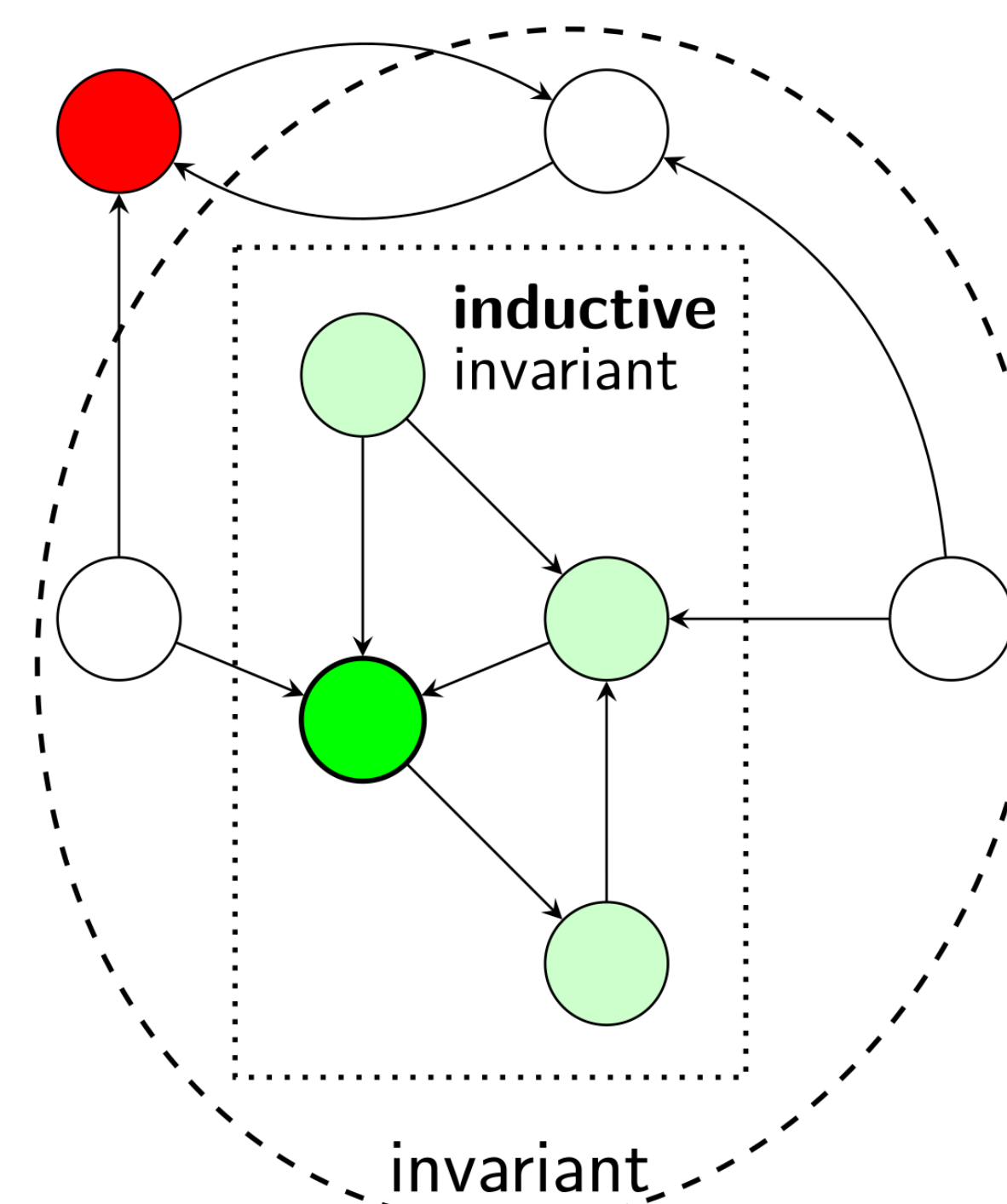
- Formal models of CPS involve **real numbers** and formal verification **requires real arithmetic**.



Formal Modelling and Verification in TLA+

Temporal Logic of Actions

- Lamport's Temporal Logic of Actions was designed to enable formal modelling and verification of concurrent systems. It enjoys excellent tool support in the form of the **TLA+ Toolbox** and has been successfully applied in industry.
- Formally proving safety specifications of discrete transition systems is typically done by finding an appropriate **invariant**.

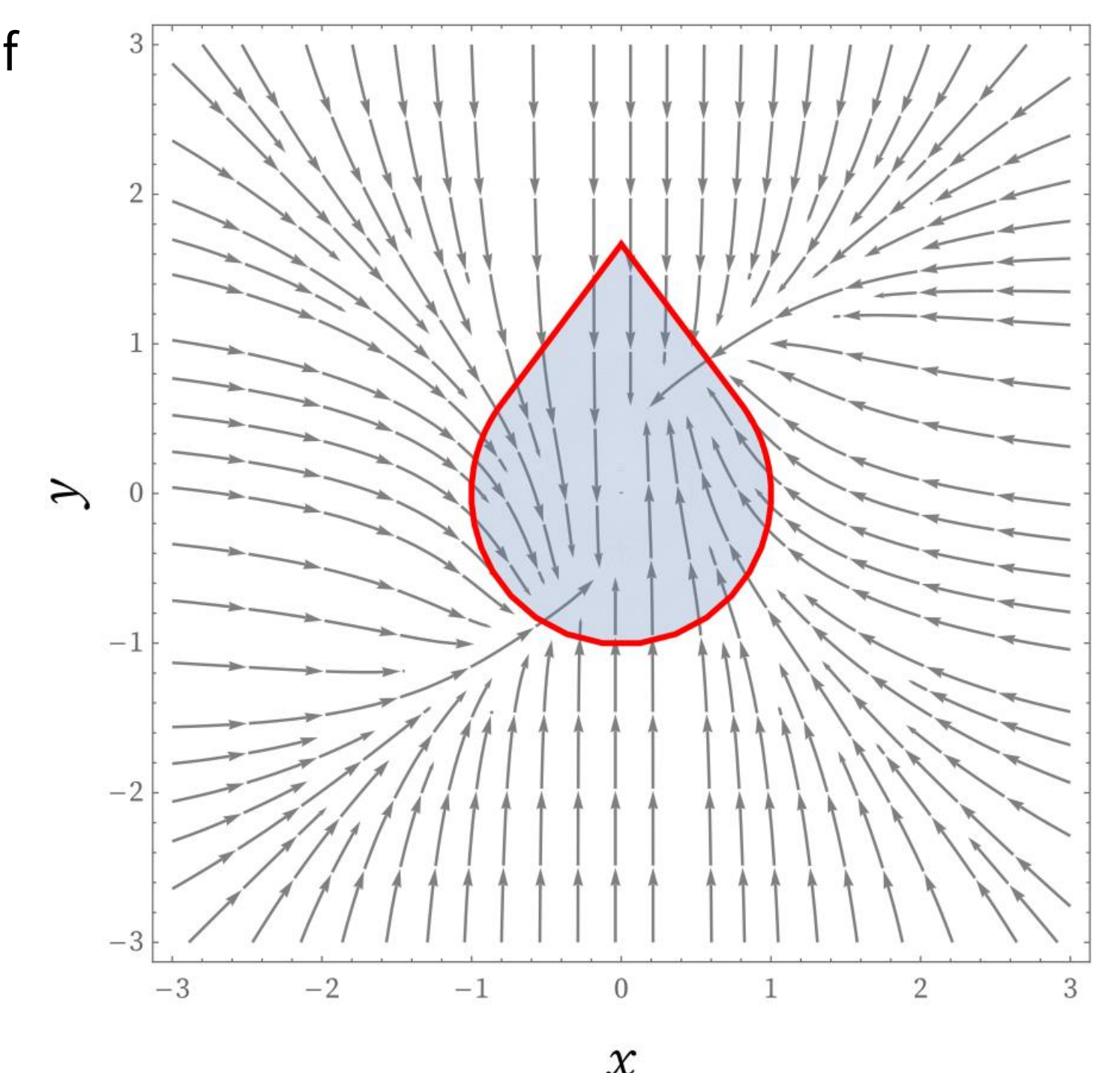


Inductive Invariants

- An **invariant** is a set of states that:
 - It includes all the initial states (as described in the safety specification).
 - It does not include any of the unsafe states.
 - The unsafe states are not reachable from the initial states.

An invariant is **inductive** if there are no transitions out of the invariant.

- A corresponding notion to an inductive invariant in continuous systems is that of a **positively invariant set / continuous invariant**.
- Recent work in computer science has established that it is **decidable** to check whether a set is positively invariant (provided it is described using polynomial functions). *This requires real arithmetic.*
- This result makes it possible to perform safety verification without having to solve the ODEs.



Automating Real Arithmetic in TLA+

- TLA+** supports real numbers (which are required for modelling and verifying CPS, especially in checking continuous invariants).
- However, the proof system currently lacks support for automatic proofs of first-order real arithmetic sentences (e.g. $\forall x, y \in \mathbb{R}. 2x^2 + (xy - y)^2 \geq -1$).
- The **TLA+ Proof Manager (TLAPM)** has now been extended to support nonlinear real arithmetic (O. V. Gunasekera et al.) – a step towards safety verification of CPS using TLA+.